

**REMARKS**

By way of the foregoing amendments to the claims, claims 1, 2, 4, and 6-9 have been amended to delete multiple dependencies and to make minor editorial changes. New claims 10-18 have been added. No new matter has been introduced by these changes.

Attached is a substitute specification which has been amended to reflect the above mentioned changes. Accompanying the substitute specification is a marked-up copy of the specification showing the changes that have been made. No new matter has been introduced by any of the changes presented in this Preliminary Amendment.

It is requested that the application be examined on the basis of the claims as amended.

Early and favorable consideration with respect to this application is respectfully requested.

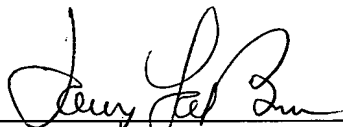
Should any questions arise in connection with this application, the undersigned respectfully requests that he be contacted at the number indicated below.

Respectfully submitted,

BUCHANAN INGERSOLL PC

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Method and associated device for generating random  
numbers in a given range

5     This disclosure is based upon French Application  
No. 0312435 filed October 24, 2003 and International  
Application No. PCT/FR2004/050510, filed October 18,  
2004, the contents of which are incorporated herein by  
reference.

10     **BACKGROUND OF THE INVENTION**

15     The invention concerns a method of obtaining a  
random number between A and B from a generator  
producing random numbers lying between 0 and W-1, with  
N the size of the numbers produced by the generator, W-  
1 the maximum value taken by the random numbers  
produced, with for example  $W = 2^N$  and A, B any integer  
numbers, less than or greater than the number W.

20     Such a situation occurs for example in an  
electronic component adapted to perform cryptographic  
calculations and comprising an N-bit random number

generator, for example  $N = 8$ . The random numbers that it can produce are thus between 0 and  $W-1 = 255$ , whilst it would be desirable to have random numbers between for example 0 and 100 or between 300 and 10000. It should be noted that it suffices to determine numbers between 0 and 9700 and then to add 300 to the number obtained in order finally to obtain a number between 300 and 10000.

Such a situation is found in practice in the majority of cryptographic applications, for example the DSA signature, the El Gamal signature or enciphering, the development of countermeasures against various attacks, etc.

Several methods are already known for producing random numbers  $R$  between 0 and  $K$  from numbers between 0 and  $W-1$ . These methods are in general implemented by software means used to control on the one hand a hardware generator that produces random numbers of size  $N$  and on the other hand calculation means performing in particular multiplication, addition, etc operations.

A first known method comprises the following steps:

a) determining the smallest integer number  $p$  such that  $K \leq WP - 1$ ,

b) producing  $p$  random numbers  $S_0, S_1, \dots, S_{p-1}$  and forming the variable  $S = \sum_{i=0}^{p-1} S_i * W^i$

c) if  $S > K$ , then returning to step b), otherwise putting  $R = S$

R is the random number sought, between 0 and K.

The equation  $S = \sum_{i=0}^{p-1} S_i * W^i$  is a representation of the variable S decomposed/recomposed in base  $(W^{p-1}, \dots, W^1, W^0)$ . It would also be possible to note  $S = S_{p-1}S_{p-2}\dots S_1S_0$ , a notation commonly used.

A second known method comprises the following steps:

a) determining the smallest integer number p such that  $K \leq W^p - 1$ ,

b) producing p random numbers  $S_0, S_1, \dots, S_{p-1}$  and forming the variable  $T = \sum_{i=0}^{p-2} S_i * W^i$  and  $S = T + S_{p-1} * W^{p-1}$

c) if  $S > K$ , putting  $R = T$  otherwise putting  $R = S$

A third known method comprises the following steps:

a) determining the smallest integer p such that  $K \leq W^p - 1$ ,

b) producing p random numbers  $S_0, S_1, \dots, S_{p-1}$  and forming the variable  $S = \sum_{i=0}^{p-1} S_i * W^i$

c) putting  $R = S \bmod (K+1)$ , that is to say the remainder of the whole-number division of S by K+1, also referred to as modular reduction of S by K+1.

These three methods can be summarised by the following steps:

a) producing  $p$  random numbers  $S_0, S_1, \dots, S_{p-1}$ ,  
 being the smallest integer number such that  $K \leq W^p - 1$   
 and forming the variable  $S = \sum_{i=0}^{p-1} S_i * W^i$

5 b) determining the random number  $R$  from the  
 variable  $S$ .

According to circumstances, during step b,  $R$  is  
 obtained from  $S$  by repeating step b (first method),  
 taking account or not of the additional random number  
 $S_{p-1}$  (second method) or performing a modular reduction  
 10 (third method).

It should be noted that, in the three methods, if  
 a number between  $A$  and  $K+A$  is required, it suffices to  
 add  $A$  to the number  $R$  obtained lying between  $0$  and  $K$ .

The main drawback of the first method is a  
 15 particularly long and especially unpredictable  
 calculation time: the step of producing the  $p$  random  
 numbers may be repeated numerous times without it being  
 possible to predict at the start the number of  
 repetitions of this step.

20 The second and third methods have the main  
 drawback of producing random numbers exhibiting a bias:  
 amongst the numbers  $R$  produced in the range  $[0, K]$ ,  
 certain values are more probable than others. In other  
 words, the numbers  $R$  produced are not perfectly random  
 25 (non-uniform distribution). This bias may have  
 significant consequences on the security of the  
 cryptographic systems liable to implement these methods.  
 The security of cryptographic systems assumes in fact  
 that the random numbers that they use are uniformly

distributed (or at least close to a uniform distribution) in the range  $[0, K]$  or  $[A, K+A]$  wished for.

Finally, the three methods are slow overall because they implement operations on large numbers, of size  $N$  (in the sense of the number of bits) greater than the size of the circuits used for the implementation. This is because the number  $K$  in particular is any number and can be greater than  $W$  and therefore of size greater than  $N$ . The variable  $S$  can also be of large size. However, the implementation of operations on large numbers requires the implementation of complex methods expensive in terms of calculation time.

### **DESCRIPTION OF THE INVENTION**

An ~~essential~~ object of the invention is to propose a method of constructing a random number  $R$  that is particularly rapid.

Thus the invention proposes a cryptographic method during which use is made of a random number generator producing random numbers  $S_i$  of size  $N$  fixed between 0 and  $W-1$ , with for example but not necessarily  $W = 2^N$ , in order to produce a random number  $R$  between 0 and a predefined limiter  $K$ .

The ~~essential~~ basic steps of a method according to the invention are as follows:

E31: a random variable  $S_i$  between 0 and  $W-1$  is produced,

E32: if the random variable  $S_i$  is strictly less than a coefficient  $K_i$  of the limiter  $K$  in base  $W$ , then the coefficient  $R_i$  of rank  $i$  of the random number  $R$  is

equal to the random variable  $S_i$  and then, for any rank  $J$  less than  $i$ , a random variable  $S_j$  between 0 and  $W-1$  is produced and  $R_j = S_j$ ,

5 E33: otherwise, if the said random variable is greater than the coefficient  $K_i$  of rank  $i$  of the limiter  $K$  in base  $W$ , then the said coefficient  $R_i$  is determined from the random variable  $S_i$  of rank  $i$  according to a predetermined function, and then the coefficient  $R_{i-1}$  is determined for the random number  $R$  of rank  $i-1$  that is  
10 immediately lower by repeating steps E31 to E33.

Thus, in a method according to the invention, the coefficients  $R_i$  of the random number  $R$  required are sought one by one, commencing with the most significant coefficient  $R_{p-1}$ . The physical generator of random  
15 numbers used thus produces random variables  $S_i$  one by one, one variable at each iteration.

In addition, the method is rapid since step E33 is executed a small number of times. This is because, as soon as one of the variables  $S_i$  produced by the physical  
20 generator is less than the associated coefficient  $K_i$  of the limiter  $K$ , the method no longer requires the processing of the variables  $S_j$  of rank less than  $i$ : thus a small number of coefficients of the number  $R$ , the most significant, are calculated the most often.

25 Finally, compared with the known methods, a method according to the invention has the advantage of working on numbers of no more than  $N$  bits,  $N$  being the size of the registers and other calculation circuits of the devices used for implementation. For example, if  $W$   
30 is equal to  $S^N$ , the coefficients  $K_i$  resulting from the

decomposition of  $K$  in base  $(W^{p-1}, \dots, W^1, W^0)$  are necessarily less than  $W$  and therefore with a size of no more than  $N$  bits. Likewise, the random variables  $S_i$  produced by the physical random number generator are also of  $N$  bits.

By adding to the ~~essential~~ basic steps an initialisation step and a step of recombination of the random number  $R$ , there are obtained:

E1: the limiter  $K$  is decomposed in base  $(W^{p-1}, W^{p-2}, \dots, W^0)$  ( $K = \sum_{i=0}^{p-1} K_i * W^i$  or  $K = K^{p-2} \dots K^1 K^0$ ),  $i$  being a loop index,  $K_i$  being a coefficient of the limiter  $K$  of rank  $i$  between 0 and  $W-1$  and  $p$  being the degree of the limiter  $K$ ,

E2: a Boolean variable  $f$  is initialised to TRUE,

E3: the following operations are performed, in a loop indexed by  $i$ ,  $i$  being an integer varying between  $p-1$  and 0:

E31: a random variable  $S_i$  between 0 and  $W-1$  is produced,

E32: if the random variable  $S_i$  is strictly less than the coefficient  $K_i$  of rank  $i$ , then the Boolean variable  $f$  is set to FALSE,

E33\_1: if the random variable  $S_i$  is strictly greater than the coefficient  $K_i$  of rank  $i$  and the Boolean variable  $f$  is TRUE, then the coefficient  $R_i$  of rank  $i$  is determined from the random variable  $S_i$  of rank  $i$  according to a predefined function,

E33\_2: otherwise  $R_i = S_i$

E34: the loop indexed  $i$  is decremented,



E4: the random number  $R$  is determined by recombination of the random coefficients  $R_i$  in base  $W$  ( $R = \sum_{i=0}^{p-1} R_i * W^i$  or  $R^{p-1} \dots R^1 R^0$ ).

5 In concrete terms, as soon as the Boolean variable  $f$  is positioned at FALSE, it remains at this value since provision is not made for repositioning it at the value TRUE, except when E2 of the method is initialised. Step E32 is executed only if the variable  $f$  is TRUE; thus, as soon as the variable  $f$  is  
10 positioned at the value FALSE, step E33\_1 is no longer executed and the method according to the invention ends rapidly.

A second objective of the invention is to propose a method of constructing random numbers whose  
15 distribution is uniform or can be made as close as desired to a uniform distribution. This objective is achieved by choosing a suitable function for the determination of the coefficient  $R_i$  from the random variable  $S_i$ .

20 According to a first embodiment of the method according to the invention, in order to determine the coefficient  $R_i$  of rank  $i$  from the random variable  $S_i$  of rank  $i$  (step E33\_1), the following substeps are performed:

25 E33\_11: if the random variable  $S_i$  is strictly greater than the coefficient  $K_i$  of the limiter  $K$ , then a new random variable  $S_i$  is produced,

E33\_12: step E33\_11 is repeated until the random variable  $S_i$  is less than the coefficient  $K_i$  of the

limiter  $K$ , and then the coefficient  $R_i$  is equalised to the random variable  $S_i$ .

5 In such an embodiment, all the coefficients  $R_i$  obtained are numbers directly produced by the hardware random number generator; and these coefficients are therefore perfect and the number  $R$  which results therefrom is also perfect. In other words the distribution obtained of the numbers  $R$  is uniform in the range  $[0, K]$ .

10 According to a second embodiment, during step E33 the coefficient  $R_i$  of rank  $i$  is chosen so as to be equal to part of the random variable  $S_i$ , a part less than the coefficient  $K_i$ . The said part corresponding in one example to a limited number of bits of the variable  $S_i$ .

15 According to a third embodiment, during step E33 the random variable  $S_i$  is reduced modulo  $K_{i+1}$ , the results of the reduction being the coefficient  $R_i$  sought.

20 These latter two embodiments are rapid compared with the known methods, essentially because the work is done on small numbers. The distributions of random numbers obtained are however not uniform: the simple fact of truncating the variable  $S_i$  or performing a reduction modulo  $K_{i+1}$  necessarily introduces a bias. 25 However, this bias is less compared with the methods of the prior art.

Moreover, it is possible to reduce the bias of the methods according to the second and third embodiments proposed, as will be seen below.

In a method according to the invention as described above, a random number  $R$  is constructed less than  $K$  from variables  $S_i$  of size  $N$  produced by a perfectly random physical generator. The number  $R$  obtained is biased, but the bias is small compared with a known method.

For this, in the second or third embodiment, a coefficient  $R_i \leq K_i$  is constructed, in particular during step E33\_1, from variables  $S_i$  of size  $N$ . In order to reduce the bias introduced on the coefficient  $R_i$ , it is proposed to construct it using the same steps E1 to E3 as for constructing the number  $R$ . In a sense, two similar methods are "interleaved". This makes it possible to reduce further the size of the numbers on which the work is carried out, and consequently to reduce further the bias on the coefficient of  $R$ , and on the final number  $R$ .

In concrete terms, in order to determine the coefficient  $R_i$  of rank  $i$  from the random variable  $S_i$  of rank  $i$  (step E33\_1), steps E1 to E4 are executed using a base  $(\beta^{q-1}, \dots, \beta^0)$  as the calculation base,  $\beta$  being an integer number strictly less than  $W$  and  $q$  being the degree of  $K_i$  in base  $\beta$ .

Step E33 is thus broken down into the following substeps:

E33\_41: the coefficient  $K_i$  of rank  $i$  of the limiter  $K$  in base  $(\beta^{q-1}, \dots, \beta^0)$  ( $K_i = \sum_{j=0}^{q-1} (K_i)_j * \beta^j$  or  $K_i = (K_i)_{q-1} \dots (K_i)_1 (K_i)_0$ ),  $j$  being a loop index,  $(K_i)_j$  being a

number between 0 and  $\beta-1$  and  $q$  being a degree of the coefficient  $K_i$ , is decomposed,

E33\_42: a second Boolean variable  $g$  is initialised to TRUE,

5 E33\_43: the following operations are performed, in a loop indexed by  $j$  varying between  $q-1$  and 0:

E33\_431: a random variable  $(S_i)_j$  between 0 and  $\beta-1$  is produced,

10 E33\_432: if the random variable  $(S_i)_j$  is strictly less than the coefficient  $(K_i)_j$ , then the second Boolean variable  $g$  is set to FALSE,

E33\_4331: if the random variable  $(S_i)_j$  is strictly greater than the coefficient  $(K_i)_j$  and the second Boolean variable  $g$  is TRUE, then a  
15 coefficient  $(R_i)_j$  is determined from the random variable  $(S_i)_j$  according to a predefined function,

E33\_4332: otherwise,  $(R_i)_j = (S_i)_j$

E33\_434: the loop indexed  $j$  is decremented,

20 E33\_44: the random number  $R_i$  is determined by recombination of the random coefficients  $(R_i)_j$  in base  $\beta$

$$(R_i = \sum_{j=0}^{q-1} (R_i)_j * \beta^j \text{ or } R_i = (R_i)_{q-1} \dots (R_i)_1 (R_i)_0).$$

As has just been seen above, by "interleaving" two methods, the bias of the random numbers  $R$  produced by the global method is reduced, whilst preserving a  
25 rapid global method. It is of course possible to imagine "interleaving" more than two methods, for example three or four, by decomposing, in step E33\_43, the numbers in base  $\gamma < \beta$ , and decomposing step E33\_43

in a succession of steps similar to steps E33\_41 to E33\_43.

5 In general terms, the more methods are "interleaved", the smaller the numbers on which the work is carried out: the duration of each step decreases and the bias of the numbers produced by the global method also decreases.

10 Another object of the invention is an electronic component adapted for implementing the method as described above. Such a component comprises in particular a generator producing random numbers of size N, and calculation circuits for performing operations on numbers of no more than N bits.

15 According to the embodiment of the method to be implemented, the calculation circuits are adapted to perform operations of comparing two numbers, number truncation and modular reduction.

20 The random number generator and the calculation circuits are preferably controlled by a software means stored in a memory of the component provided for this purpose.

The invention also concerns a chip card comprising an electronic component as described above.

## CLAIMS

1. A cryptographic method during which use is made of a random number generator producing random numbers  $S_i$  of size  $N$  fixed between 0 and  $W-1$ , in order to produce a random number  $R$  between 0 and a predefined limiter  $K$ , characterised in that:

E31: a random variable  $S_i$  between 0 and  $W-1$  is produced,

E32: if the random variable  $S_i$  is strictly less than a coefficient  $K_i$  of the limiter  $K$  in base  $W$ , then the coefficient  $R_i$  of rank  $i$  of the random number  $R$  is equal to the random variable  $S_i$  and then, for any rank  $J$  less than  $i$ , a random variable  $S_j$  between 0 and  $W-1$  is produced and  $R_j = S_j$ ,

E33: otherwise, if the said random variable is greater than the coefficient  $K_i$  of rank  $i$  of the limiter  $K$  in base  $W$ , then the said coefficient  $R_i$  is determined from the random variable  $S_i$  of rank  $i$  according to a predetermined function, and then the coefficient  $R_{i-1}$  is determined for the random number  $R$  of rank  $i-1$  that is immediately lower by repeating steps E31 to E33.

2. A method according to claim 2, during which the following steps are performed:

E1: the limiter  $K$  is decomposed in base  $(W^{p-1}, W^{p-2}, \dots, W^0)$  in the form  $K = \sum_{i=0}^{p-1} K_i * W^i$ ,  $i$  being a loop index,  $K_i$  being a coefficient of the limiter  $K$  of rank  $i$  between 0 and  $W-1$  and  $p$  being the degree of the limiter  $K$ ,

E2: a Boolean variable  $f$  is initialised to TRUE,  
 E3: the following operations are performed, in a loop indexed by  $i$ ,  $i$  being an integer varying between  $p-1$  and 0:

5           E31: a random variable  $S_i$  between 0 and  $W0-1$  is produced,

          E32: if the random variable  $S_i$  is strictly less than the coefficient  $K_i$  of rank  $i$ , then the Boolean variable  $f$  is set to FALSE,

10           E33\_1: if the random variable  $S_i$  is strictly greater than the coefficient  $K_i$  of rank  $i$  and the Boolean variable  $f$  is TRUE, then the coefficient  $R_i$  of rank  $i$  is determined from the random variable  $S_i$  of rank  $i$  according to a predefined function,

15           E33\_2: otherwise  $R_i = S_i$

          E34: the loop indexed  $i$  is decremented,

E4: the random number  $R$  is determined by recombination of the random coefficients  $R_i$  in base  $W$  according to the equation:  $R = \sum_{i=0}^{p-1} R_i * W^i$ .

20           3. A method according to claim 2, during which, in order to determine the coefficient  $R_i$  of rank  $i$  from the random variable  $S_i$  of rank  $i$  (steps E33\_1 and E33\_2), the following substeps are performed:

          E33\_11: if the random variable  $S_i$  is strictly greater than the coefficient  $K_i$  of the limiter  $K$ , then a new random variable  $S_i$  is produced,

25           E33\_12: step E33\_11 is repeated until the random variable  $S_i$  is less than the coefficient  $K_i$  of the

limiter  $K$ , and then the coefficient  $R_i$  is equalised to the random variable  $S_i$ .

4. A method according to claim 2, during which the coefficient  $R_i$  of rank  $i$  is chosen (steps E33-1 and E33\_2) equal to the part of the random variable  $S_i$ , the part less than the coefficient  $K_i$ , the said part corresponding to a limited number of bits of the variable  $S_i$ .

5. A method according to claim 2, during which, in order to determine the coefficient  $R_i$  of rank  $i$  from the random variable  $S_i$  of rank  $i$  (step E33), the random variable  $S_i$  is reduced modulo  $K_i+1$ , the result of the reduction being the coefficient sought.

6. A method according to one of claims 1 to 5, during which, in order to determine the coefficient  $R_i$  of rank  $i$  from the random variable  $S_i$  of rank  $i$  (step E33), steps E1 to E4 are executed using a base  $(\beta^{q-1}, \dots, \beta^0)$  as the calculation base,  $\beta$  being an integer strictly less than  $W$  and  $q$  being the degree of  $k$  in case  $\beta$ .

7. A method according to claim 6, in which step E33 is broken down into the following substeps:

E33\_41: the coefficient  $K_i$  of rank  $i$  of the limiter  $K$  in base  $(\beta^{q-1}, \dots, \beta^0)$  in the form

$$K_i = \sum_{j=0}^{q-1} (K_i)_j * \beta^j, \quad j \text{ being a loop index, } (K_i)_j \text{ being a}$$

number between 0 and  $\beta-1$  and  $q$  being a degree of the coefficient  $K_i$ , is decomposed,



E33\_42: a second Boolean variable  $g$  is initialised to TRUE,

E33\_43: the following operations are performed, in a loop indexed by  $j$  varying between  $q-1$  and  $0$ :

5       E33\_431: a random variable  $(S_i)_j$  between  $0$  and  $\beta-1$  is produced,

E33\_432: if the random variable  $(S_i)_j$  is strictly less than the coefficient  $(K_i)_j$ , then the second Boolean variable  $g$  is set to FALSE,

10       E33\_4331: if the random variable  $(S_i)_j$  is strictly greater than the coefficient  $(K_i)_j$  and the second Boolean variable  $g$  is TRUE, then a coefficient  $(R_i)_j$  is determined from the random variable  $(S_i)_j$  according to a predefined function,

15       E33\_4332: otherwise,  $(R_i)_j = (S_i)_j$

E33\_434: the loop indexed  $j$  is decremented,

E33\_44: the random number  $R_i$  is determined by recombination of the random coefficients  $(R_i)_j$  in base  $\beta$  according to the equation:  $R_i = \sum_{j=0}^{q-1} (R_i)_j * \beta^j$ .

20       8. An electronic component comprising a generator of random numbers of size  $N$ , calculation circuits performing in particular a comparison, a truncation and/or a modular reduction on numbers of no more than  $N$  bits, and a means of controlling the random  
25       number generator and calculation circuits, the said control means being adapted for implementing a method according to one of claims 1 to 7.

9. A chip card comprising an electronic component according to the preceding claim.

ABSTRACT

The invention relates to a cryptographic method wherein  
a random number generator producing random numbers  $S_i$   
5 whose size  $N$  is fixed between 0 and  $W-1$  is used to  
produce a random number  $R$  between 0 and a predefined  
limiter  $K$ . According to the invention: E31: a random  
variable  $S_i$  is produced, ranging from 0- $W-1$ , E32: if the  
random variable  $S_i$  is strictly lower than a coefficient  
10  $K_i$  of the limiter  $K$  in base  $W$ , the coefficient  $R_i$  of  
order  $i$  of the random number  $R$  is equal to the random  
number  $S_i$  then, for all orders  $j$  which are lower than  $i$ ,  
a random variable  $S_j$  of 0- $W-1$  is produced and  $R_j=S_j$ .  
E33: unless, if said random variable is greater than  
15 coefficient  $K_i$  of position  $i$  of the limiter  $K$  is base  $W$ ,  
whereupon said coefficient  $R_i$  is determined on the basis  
of the random variable  $S_i$  of order  $i$  according to a  
predetermined function, then a coefficient  $R_{i-1}$  is  
determined for the random number  $R$  of order  $i-1$  which  
20 is immediately lower by repeating stages E31 - E33.  
The invention also relates to an electronic component  
which is adapted for implementation of said method and  
a chip card with said component integrated therein.  
The invention can be applied to cryptographic  
25 calculation.